



Cryptography of a Gray Level Image using a Novel Block Cipher Involving Feistel Structure and Modular Arithmetic

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Abstract: In this paper, we have made use of a novel block cipher involving Feistel structure and modular arithmetic for encrypting a Gray level image. The image is represented in the form of a matrix of size 256×256 , and this is divided into 32 parts, wherein each part is of size 32×64 . A key of size 8×8 is taken and it is transformed into a key of size 32×32 . This key is utilised in carrying out the computations required for encryption and decryption. As the cipher is a strong one, we find that no one can identify the original image by any means.

Keywords: Block Cipher, Gray level image, Feistel structure, Modular arithmetic, Encryption, Decryption and Encrypted image.

I. INTRODUCTION

The study of cryptography of images has been an interesting area of research in recent years. Most of the symmetric block ciphers [1–4] and the public key ciphers [5–7] are utilised in the development of the image cryptography. In all these investigations, security of images is achieved in a significant manner.

In a recent paper [8], we have developed a novel block cipher by using Feistel structure and modular arithmetic. In this analysis, we have made use of '+' operation instead of XOR operation, which was used in the classical Feistel cipher [9]. Here, we have shown that the cipher is a strong one as the length of the key is quite considerable and it leads to nonlinearities as we have included the key on both the sides of a portion of the plain text.

In the present paper, our objective is to study the encryption and decryption of a gray level image. To this end, we have followed the procedure which is the same as that utilised in [10].

In Section 2, we have presented the development of the procedure for the cryptography of a gray level image. In Section 3, we have given an example and illustrated the process. Finally, in Section 4, we have mentioned the computations carried out in this analysis and drawn conclusions.

II. DEVELOPMENT OF A PROCEDURE FOR THE CRYPTOGRAPHY OF A GRAY LEVEL IMAGE

Let us consider a plain text P having $2m^2$ characters. This can be written in the form of a pair of square matrices denoted

as P_0 and Q_0 , in which each is having m^2 characters. Let K be a key (square) matrix of size m .

Now, the process of encryption and the process of decryption are governed by the relations

$$\begin{aligned} P_i &= Q_{i-1}, \\ Q_i &= (P_{i-1} + (F(Q_{i-1}, K)) \bmod N), \text{ for } i = 1 \text{ to } n, \end{aligned}$$

$$F(Q_{i-1}, K) = (KQ_{i-1}K) \bmod N, \quad (2.1)$$

and

$$\begin{aligned} Q_{i-1} &= P_i, \\ P_{i-1} &= (Q_i + (F(P_i, K)) \bmod N), \text{ for } i = n \text{ to } 1, \\ F(P_i, K) &= (KP_iK) \bmod N. \end{aligned} \quad (2.2)$$

In the present analysis, N will be chosen appropriately. Here, n stands for the number of rounds of the iteration process.

The flow charts describing encryption and decryption are shown in Fig. 1.

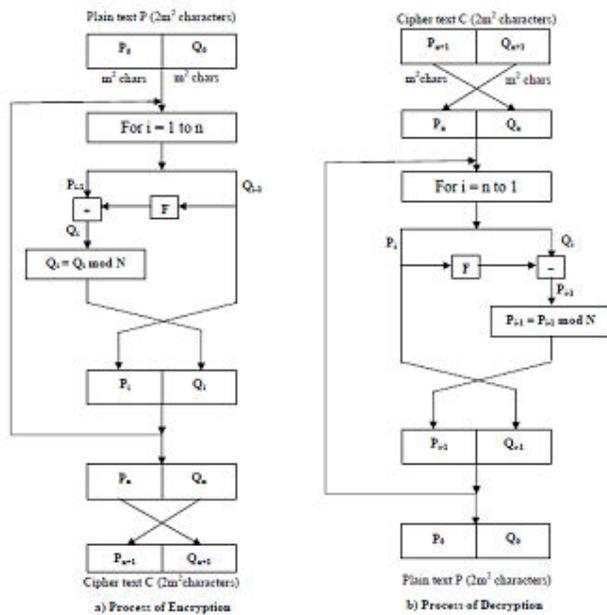


Figure 1. Schematic diagram of cipher

Here, we present the algorithms for encryption and decryption.

Algorithm for Encryption

1. Read P, m, n, N
2. $P_0 =$ Left half of P
 $Q_0 =$ Right half of P
3. for $i = 1$ to n
{
 $P_i = Q_{i-1}$
 $F = (KQ_{i-1}K) \bmod N$
 $Q_i = (P_{i-1} + F) \bmod N$
}
4. $P_{n+1} = Q_n$
 $Q_{n+1} = P_n$
5. $C = P_{n+1} \parallel Q_{n+1}$ /* \parallel stands for concatenation
6. Write (C)

Algorithm for Decryption

1. Read C, m, n, N
2. $P_{n+1} =$ Left half of C
 $Q_{n+1} =$ Right half of C
3. for $i = n$ to 1
{
 $Q_{i-1} = P_i$
 $F = (KP_iK) \bmod N$
 $P_{i-1} = (Q_i - F) \bmod N$
}
4. $P_0 = Q_1$
 $Q_0 = P_1$
5. $P = P_0 \parallel Q_0$ /* \parallel stands for concatenation
6. Write (P)

III. ILLUSTRATION OF THE CRYPTOGRAPHY OF AN IMAGE

Let us consider a small sample gray level image. Let this be represented in the form of a matrix containing 8 rows and 16 columns. This is given by

$$P = \begin{pmatrix} 001 & 001 & 004 & 008 & 000 & 006 & 011 & 002 & 007 & 004 & 002 & 003 & 003 & 003 & 005 & 009 \\ 001 & 002 & 001 & 003 & 012 & 000 & 001 & 011 & 003 & 003 & 010 & 004 & 002 & 006 & 012 & 010 \\ 001 & 010 & 004 & 001 & 001 & 001 & 001 & 006 & 000 & 000 & 011 & 000 & 012 & 117 & 163 & 147 \\ 001 & 003 & 005 & 005 & 001 & 011 & 019 & 001 & 001 & 012 & 001 & 104 & 214 & 211 & 205 & 239 \\ 001 & 010 & 001 & 014 & 016 & 001 & 001 & 001 & 007 & 081 & 197 & 252 & 250 & 226 & 236 & 240 \\ 001 & 001 & 015 & 001 & 001 & 014 & 002 & 010 & 153 & 253 & 251 & 248 & 245 & 239 & 238 & 241 \\ 017 & 001 & 001 & 022 & 001 & 067 & 230 & 253 & 249 & 251 & 240 & 223 & 248 & 239 & 214 & 233 \\ 001 & 026 & 000 & 000 & 005 & 107 & 253 & 242 & 250 & 229 & 240 & 248 & 234 & 241 & 248 & 229 \end{pmatrix} \quad (3.1)$$

Let us take a key matrix K of size 8 x 8 in the form

$$K = \begin{pmatrix} 175 & 173 & 027 & 065 & 032 & 065 & 017 & 076 \\ 232 & 084 & 072 & 069 & 032 & 185 & 069 & 082 \\ 027 & 179 & 102 & 033 & 083 & 097 & 073 & 032 \\ 065 & 084 & 143 & 069 & 105 & 153 & 213 & 163 \\ 184 & 028 & 049 & 005 & 069 & 031 & 166 & 109 \\ 208 & 185 & 077 & 234 & 207 & 171 & 071 & 080 \\ 237 & 249 & 101 & 057 & 095 & 191 & 037 & 132 \\ 127 & 107 & 032 & 085 & 117 & 254 & 165 & 087 \end{pmatrix} \quad (3.2)$$

On applying the encryption algorithm given in section 2, we have

$$C = \begin{pmatrix} 176 & 322 & 268 & 258 & 278 & 346 & 134 & 278 & 172 & 302 & 280 & 138 & 222 & 134 & 254 & 182 \\ 264 & 238 & 210 & 180 & 150 & 116 & 400 & 304 & 398 & 398 & 212 & 234 & 292 & 160 & 286 & 090 \\ 240 & 122 & 318 & 092 & 446 & 112 & 132 & 068 & 334 & 418 & 312 & 290 & 340 & 140 & 106 & 342 \\ 348 & 462 & 344 & 134 & 138 & 450 & 210 & 366 & 204 & 218 & 320 & 402 & 174 & 270 & 486 & 402 \\ 140 & 096 & 282 & 112 & 100 & 370 & 210 & 056 & 034 & 172 & 252 & 254 & 238 & 106 & 156 & 034 \\ 024 & 164 & 420 & 226 & 212 & 168 & 180 & 278 & 208 & 180 & 436 & 446 & 468 & 216 & 340 & 370 \\ 460 & 276 & 366 & 206 & 048 & 474 & 286 & 218 & 356 & 076 & 270 & 394 & 296 & 274 & 102 & 118 \\ 044 & 330 & 160 & 362 & 242 & 338 & 170 & 326 & 368 & 102 & 356 & 458 & 370 & 254 & 234 & 254 \end{pmatrix} \quad (3.3)$$

This can be brought to the form of an image (Encrypted image) given in Fig. 2.

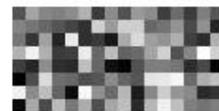


Figure 2. Encrypted form of the sample Image

On applying the decryption algorithm (Section 2) on the cipher text in (3.3), we get back the original plain text P.

The procedure described above can be adapted to any gray level image of any size by taking a suitable key and dividing the image into a set of subimages appropriately. In the next section, we describe the encryption and decryption processes of a real image.

IV. COMPUTATIONS AND CONCLUSIONS

Let us consider the image of Dr. S. Radhakrishnan, given in Fig. 3.



Figure 3 Image of Dr. S. Radhakrishnan

This gray level image is represented in the form of a matrix of size 256 x 256. In order to carryout the process of encryption in a convenient manner, this image is divided into 32 parts, in which each part is of size 32 x 64.

We consider a key matrix of size 32 x 32, which is generated from the key matrix K of size 8 x 8 given in (3.2), by applying the procedure described in [10].

On using the key K of size 32 x 32, and the procedure mentioned in section 3, we have encrypted all the 32 parts of the image. Thus, we have found the cipher text corresponding to the entire image. This is exhibited in Fig. 4.

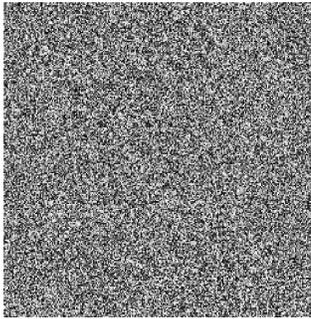


Figure 4. Encrypted form of the Entire image

On carrying out the process of decryption on all the 32 parts separately, we have got the original plain text and the corresponding image (Fig. 3).

All the computations in this analysis are carried out by writing C programs for the encryption and the decryption algorithms (given in section 2). We have used MATLAB in the development of encrypted image.

As the image corresponding to the cipher text is totally in a peculiar form, we conclude that no one can recognize it in the process of transmission. This is all due to the strength of the cipher.

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