



TOPOLOGICAL PROPERTIES OF NETWORK AND INFORMATION FLOW FOR PARALLEL AND DISTRIBUTED SYSTEM

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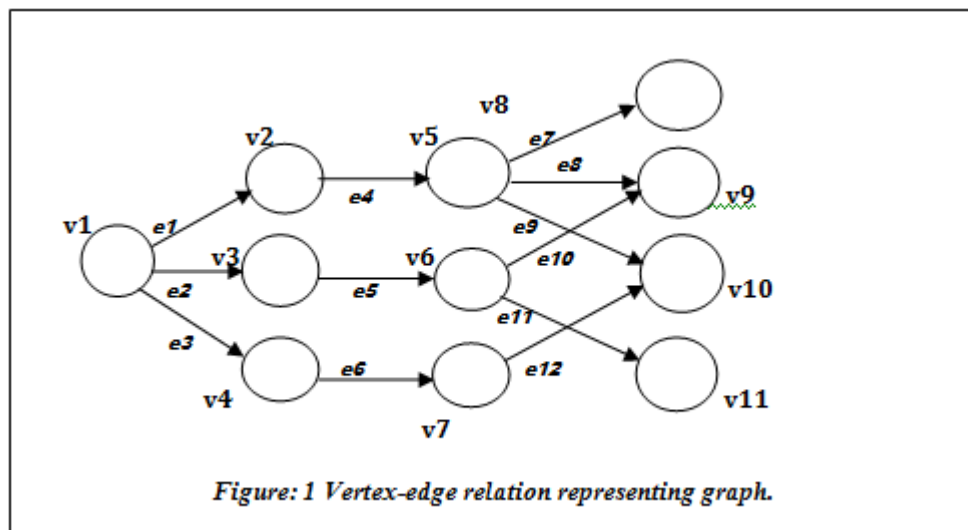
Abstract: In this paper we are analyzing and developing algorithms for the connectivity and complexity of point to point network. For the development of communication network algorithms we will convert the point to point network into its connectivity matrix. These connectivity matrix relations will be depending on the Processors to Processors relation or processor to edge relation or edge to edge relations. The logical operations will be performed between the vectors of nodes or edges of the connectivity matrix. These operations will be used to explore the topological properties of the architecture for information flow.

Keywords: Interconnection Network, Point to Point network, Connectivity matrix algorithms.

INTRODUCTION

Let $G = (v, e)$ is a graph where v is the set of vertexes of a point-to-point communication network and e is the set of edges. The information

can be sent noiselessly from node i to node j (for all $(i, j) \in e$) through edges of this communication network.



Now we are analyzing the connectivity and complexity of point to point network by converting the network into its equivalent connectivity matrix so that communication algorithm to connectivity can be developed. Here

we will use vectors of the matrix for analyzing the connectivity of nodes. The logical operation between vectors will be used for the study of topological property of the architecture for information flow [1] [2].

	<i>e1</i>	<i>e2</i>	<i>e3</i>	<i>e4</i>	<i>e5</i>	<i>e6</i>	<i>e7</i>	<i>e8</i>	<i>e9</i>	<i>e10</i>	<i>e11</i>	<i>e12</i>	Degree
<i>v1</i>	1	1	1	0	0	0	0	0	0	0	0	0	3
<i>v2</i>	1	0	0	1	0	0	0	0	0	0	0	0	2
<i>v3</i>	0	1	0	0	1	0	0	0	0	0	0	0	2
<i>v4</i>	0	0	1	0	0	1	0	0	0	0	0	0	2
<i>v5</i>	0	0	0	1	0	0	1	1	0	1	0	0	4
<i>v6</i>	0	0	0	0	1	0	0	0	1	0	1	0	3
<i>v7</i>	0	0	0	0	0	1	0	0	0	0	0	1	2
<i>v8</i>	0	0	0	0	0	0	1	0	0	0	0	0	4
<i>v9</i>	0	0	0	0	0	0	0	1	1	0	0	0	3
<i>v10</i>	0	0	0	0	0	0	0	0	0	1	0	1	4
<i>v11</i>	0	0	0	0	0	0	0	0	0	0	1	0	3
No of Vertices (Edges value)	2	2	2	2	2	2	2	2	2	2	2	2	

v1v2 *v1v3* *v1v4* *v2v5* *v3v6* *v4v7* *v5v8* *v5v9* *v6v9* *v5v10* *v6v10* *v7v10*

Table.1 (a) vertex to edges relation of figure 1.

STRUCTURAL RELATION OF NODES AND VECTORS

Here we are representing the structural relation between nodes & edges of any graph by converting matrices. We are assuming that the relation between node to itself or edge to itself is possible and we will take 1 for self node/self edge in our connectivity matrix.

The table 1 gives us degree of each node and number of vertexes connected with each edge. The table 2 shows vertex to vertex relation of figure 1. In table 3 we have represented edge to edge & vertex to vertex relations of figure 2 and figure 3. Table 1(b) shows the logical operation of edges with its compliments to shows tautology or contradiction for the purpose to validate the connectivity.

		Logical operation of edge & compliment of each edge	
edges	Compliment of each edge	with "AND"	with "OR"
<i>e1</i> (11000000000)	<i>e1</i> (00111111111)	(00000000000)	(11111111111)
<i>e2</i> (10100000000)	<i>e2</i> (01011111111)	(00000000000)	(11111111111)
<i>e3</i> (10010000000)	<i>e3</i> (01101111111)	(00000000000)	(11111111111)
<i>e4</i> (01001000000)	<i>e4</i> (10110111111)	(00000000000)	(11111111111)
<i>e5</i> (00100100000)	<i>e5</i> (11010011111)	(00000000000)	(11111111111)
<i>e6</i> (00010010000)	<i>e6</i> (11101101111)	(00000000000)	(11111111111)
<i>e7</i> (00001001000)	<i>e7</i> (11110110111)	(00000000000)	(11111111111)
<i>e8</i> (00001000100)	<i>e8</i> (11110110111)	(00000000000)	(11111111111)
<i>e9</i> (00000100100)	<i>e9</i> (11111011011)	(00000000000)	(11111111111)
<i>e10</i> (000001000010)	<i>e10</i> (11110111101)	(00000000000)	(11111111111)
<i>e11</i> (00000100001)	<i>e11</i> (11111011110)	(00000000000)	(11111111111)
<i>e12</i> (00000010010)	<i>e12</i> (11111101101)	(00000000000)	(11111111111)

Table 1(b) Logical Operation of edge with its compliments.

	v1	v2	v3	v4	v5	v6	v7	v8	v9	v10	v11	Degree
v1	0	1	1	1	0	0	0	0	0	0	0	3
v2	0	0	0	0	1	0	0	0	0	0	0	1
v3	0	0	0	0	0	1	0	0	0	0	0	1
v4	0	0	0	0	0	0	1	0	0	0	0	1
v5	0	0	0	0	0	0	0	1	1	1	0	3
v6	0	0	0	0	0	0	0	0	1	0	1	2
v7	0	0	0	0	0	0	0	0	0	1	0	1
v8	0	0	0	0	0	0	0	0	0	0	0	0
v9	0	0	0	0	0	0	0	0	0	0	0	0
v10	0	0	0	0	0	0	0	0	0	0	0	0
v11	0	0	0	0	0	0	0	0	0	0	0	0

Table. 2 vertex to vertex relation of figure 1.

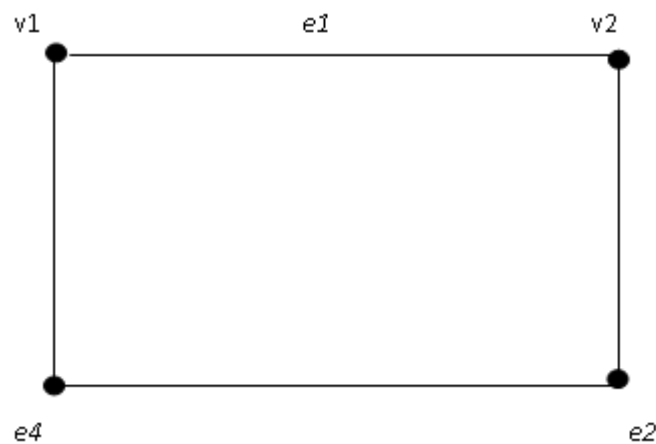


Fig.2 Graph Representation

	e1	e2	e3	e4
e1	1	1	0	1
e2	1	1	1	0
e3	0	1	1	1
e4	1	0	1	1

	v1	v2	v3	v4
v1	1	1	0	1
v2	1	1	1	0
v3	0	1	1	1
v4	1	0	1	1

Table.3 edge - edge & vertex – vertex Operation.

Table 5 shows the logical relation between processors of figure 3. In figure 3 we have two graph which are converted in vertex – vertex relations as tables shown in table3. Now we are

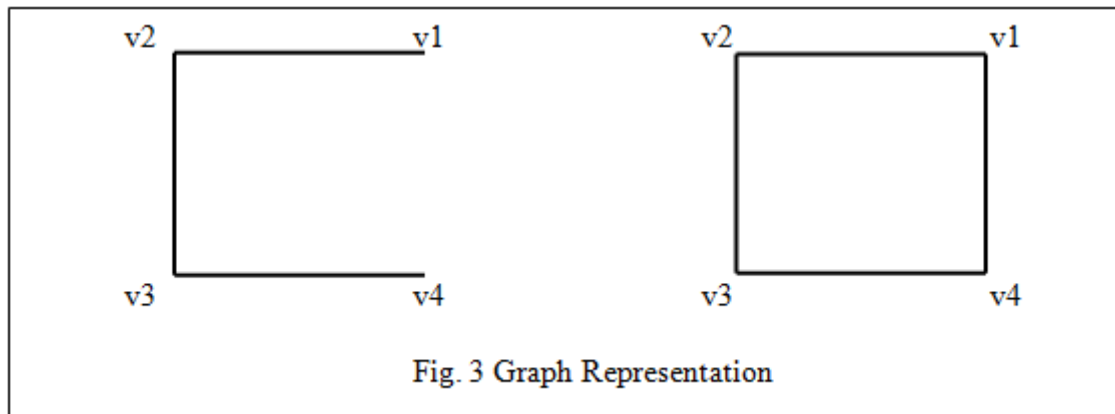
using logical operation "AND" between vectors of matrix shown in table3. We have shown that "AND" gives common connectivity between two nodes.

$$v_i \text{ AND } v_j \quad \left\{ \begin{array}{l} = 1 \text{ where } i=j=1 \\ 0 \text{ other wise} \end{array} \right.$$

The "AND" operation between compliments of nodes will give us contradiction where as "OR"

operation between complimentary nodes gives tautology[3].

$$v_i \text{ OR } v_j \quad \left\{ \begin{array}{l} 0 \text{ if } i=j=1 \\ 1 \text{ other wise} \end{array} \right.$$



	v1	v2	v3	v4
v1	0	1	0	0
v2	1	0	1	0
v3	0	1	0	1
v4	0	0	1	0

(a)

	v1	v2	v3	v4
v1	0	1	0	1
v2	1	0	1	1
v3	0	1	0	1
v4	1	0	1	0

(b)

Table.4 (a) & (b) Matrix for vertex to vertex Operation.

Possible Connectivity of nodes with AND Operation:

Nodes	(a) AND	Possible Connectivity of nodes	(b)AND	Possible Connectivity of nodes
v1,v2	$v1(0\ 1\ 0\ 0) \wedge v2(1\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil	$v1(0\ 1\ 0\ 1) \wedge v2(1\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil
v1,v3	$v1(0\ 1\ 0\ 0) \wedge v3(0\ 1\ 0\ 1) = \{0\ 1\ 0\ 0\}$	(2)=1	$v1(0\ 1\ 0\ 1) \wedge v3(0\ 1\ 0\ 1) = \{0\ 1\ 0\ 1\}$	(2,4)=2
v1,v4	$v1(0\ 1\ 0\ 0) \wedge v4(0\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil	$v1(0\ 1\ 0\ 1) \wedge v4(1\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil
v2,v3	$v2(1\ 0\ 1\ 0) \wedge v3(0\ 1\ 0\ 1) = \{0\ 0\ 0\ 0\}$	Nil	$v2(1\ 0\ 1\ 0) \wedge v3(0\ 1\ 0\ 1) = \{0\ 0\ 0\ 0\}$	Nil
v2,v4	$v2(1\ 0\ 1\ 0) \wedge v4(0\ 0\ 1\ 0) = \{0\ 0\ 1\ 0\}$	(3)=1	$v2(1\ 0\ 1\ 0) \wedge v4(1\ 0\ 1\ 0) = \{1\ 0\ 1\ 0\}$	(1,3)=2
v3,v4	$v3(0\ 1\ 0\ 1) \wedge v4(0\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil	$v3(0\ 1\ 0\ 1) \wedge v4(1\ 0\ 1\ 0) = \{0\ 0\ 0\ 0\}$	Nil

Table 5 Matrix relation with AND operation.

Possible Connectivity of Nodes with OR operation

Nodes	(a) OR	Possible Connectivity of nodes	(b)OR	Possible Connectivity of nodes
$v1, v2$	$v1(0\ 1\ 0\ 0) \vee v2(1\ 0\ 1\ 0) = \{1\ 1\ 1\ 0\}$	$(1,2,3)=3$	$v1(0\ 1\ 0\ 1) \vee v2(1\ 0\ 1\ 0) = \{1\ 1\ 1\ 1\}$	$(1,2,3,4)=4$
$v1, v3$	$v1(0\ 1\ 0\ 0) \vee v3(0\ 1\ 0\ 1) = \{0\ 1\ 0\ 1\}$	$(2,4)=2$	$v1(0\ 1\ 0\ 1) \vee v3(0\ 1\ 0\ 1) = \{0\ 1\ 0\ 1\}$	$(2,4)=2$
$v1, v4$	$v1(0\ 1\ 0\ 0) \vee v4(0\ 0\ 1\ 0) = \{0\ 1\ 1\ 0\}$	$(2,3)=2$	$v1(0\ 1\ 0\ 1) \vee v4(1\ 0\ 1\ 0) = \{1\ 1\ 1\ 1\}$	$(1,2,3,4)=4$
$v2, v3$	$v2(1\ 0\ 1\ 0) \vee v3(0\ 1\ 0\ 1) = \{1\ 1\ 1\ 1\}$	$(1,2,3,4)=4$	$v2(1\ 0\ 1\ 0) \vee v3(0\ 1\ 0\ 1) = \{1\ 1\ 1\ 1\}$	$(1,2,3,4)=4$
$v2, v4$	$v2(1\ 0\ 1\ 0) \vee v4(0\ 0\ 1\ 0) = \{1\ 0\ 1\ 0\}$	$(1,3)=2$	$v2(1\ 0\ 1\ 0) \vee v4(1\ 0\ 1\ 0) = \{1\ 0\ 1\ 0\}$	$(1,3)=2$
$v3, v4$	$v3(0\ 1\ 0\ 1) \vee v4(0\ 0\ 1\ 0) = \{0\ 1\ 1\ 1\}$	$(2,3,4)=3$	$v3(0\ 1\ 0\ 1) \vee v4(1\ 0\ 1\ 0) = \{1\ 1\ 1\ 1\}$	$(1,2,3,4)=4$

Table 6 Matrix relation with OR operation.

CONCLUSION AND FUTUREWORK

In this paper we are trying to establish the connectivity and complexity of nodes for the development of algorithms. The following are the conclusion of vector operations between nodes of connectivity matrix.

1. The logical operation "AND" shows the connectivity of nodes in "Race" condition.
2. The logical operation "OR" shows the connectivity of nodes when any of the node is active.
3. The complimentary nodes with "AND" not possible because of contradiction where as complements nodes with is always possible because of tautology.
4. The use of logical operations shows the validity of possible communication between nodes. The algorithms will

validate the communication is possible or not at some point. We will show that binary relations are always possible in our future studies.

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